

Pure Exploration by Solving Games

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Main recipe

Take two adversarial strategies for regret minimization.

Add optimism.

Get one stochastic bandit algorithm for pure exploration.

Paper: Rémy Degenne, Wouter M. Koolen and Pierre Ménard, Non-Asymptotic Pure Exploration by Solving Games, NeurIPS 2019.

Pure Exploration

Usual Queries

- Best Arm Identification
- Thresholding Bandit

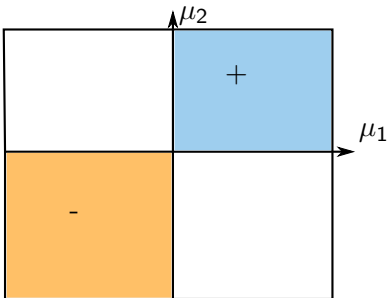
Our setting

- Bandit parametrized by means $\mu \in \mathcal{M} \subset \mathbb{R}^K$.
- Answers \mathcal{I} . Correct answer function $i^* : \mathcal{M} \rightarrow \mathcal{I}$.
- Fixed confidence $\delta \in [0, 1]$.
- Algorithm stops at time τ_δ , returns \hat{i} .

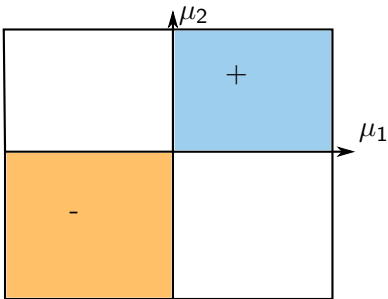
Goal: δ -correct algorithm, such that

$$\forall \mu \in \mathcal{M} \quad \mathbb{P}_\mu(\hat{i} \neq i^*(\mu)) \leq \delta, \quad \mathbb{E}_\mu \tau_\delta \text{ is minimal.}$$

Pure Exploration



Pure Exploration



This talk: about sampling rules.

Use GLRT stopping rule from Garivier and Kaufmann, 2016.

Sample complexity: what is
“minimal”?

Lower Bound

Any δ -correct algorithm on \mathcal{M} must verify for all $\mu \in \mathcal{M}$,

$$\mathbb{E}_{\mu}[\tau_{\delta}] \max_{w \in \Delta_K} \inf_{\lambda \in \neg i^*(\mu)} \sum_{k=1}^K w^k d(\mu^k, \lambda^k) \geq \text{kl}(\delta, 1 - \delta)$$

$$\neg i = \{\lambda \in \mathcal{M} : i^*(\lambda) \neq i\}.$$

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Follow the lower bound: attempt 1

Track and Stop

Compute estimated problem $\hat{\mu}_t$.Compute the solution w_t^* to

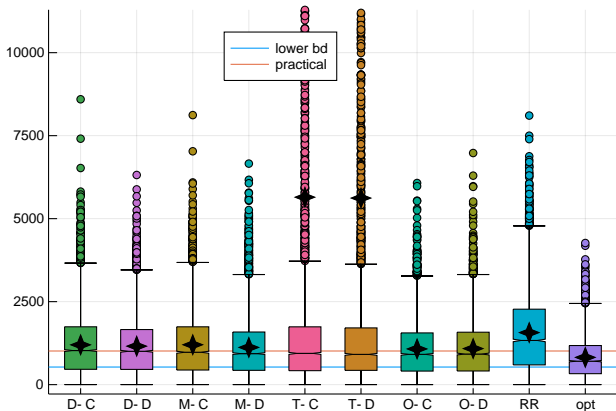
$$\operatorname{argmax}_{w \in \Delta_K} \inf_{\lambda \in \neg j^*(\hat{\mu}_t)} \sum_{k=1}^K w^k d(\hat{\mu}_t^k, \lambda^k).$$

If an arm is sampled less than \sqrt{t} , sample it (forced exploration).Otherwise, sample arm $k_t = \operatorname{argmin} N_{t-1}^k - (w_t^*)^k$ (tracking).

Track-and-Stop

- Asymptotically optimal,
- But sometimes only asymptotically.

$$\liminf_{\delta \rightarrow 0} \frac{\mathbb{E}_{\mu} \tau_{\delta}}{\log(1/\delta)} \leq \frac{1}{\sup_{w \in \Delta_K} \inf_{\lambda \in \neg i^*(\mu)} \sum_{k=1}^K w^k d(\mu^k, \lambda^k)} .$$



Follow the lower bound: attempt 2

with games!

A Game

Suppose μ , $i = i^*(\mu)$ known.

- k-Player plays in $\{1, \dots, K\}$.
- λ -Player plays in $\neg i$.
- zero-sum. reward for k-player: $d(\mu^k, \lambda^k)$.

After t iterations: reward $\sum_{s=1}^t d(\mu^{k_s}, \lambda_s^{k_s})$.

Algorithms

- Regret-minimizing algorithm for k: AdaHedge.
- Regret-minimizing algorithm for λ : Best-Response.
- Result: value $\frac{1}{t} \sum_{s=1}^t d(\mu^{k_s}, \lambda_s^{k_s})$ converges to max-min.

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Algorithm for Pure Exploration

At stage $t \in \mathbb{N}$,

- Compute $\hat{\mu}_t$, define candidate answer i_t .
- Define game with **optimistic reward** $\max_{\xi \in [\hat{\mu}_t^k \pm \dots]} d(\xi, \lambda^k)$.
- Do 1 iteration of each learner on optimistic game.
- Sample the arm prescribed by the k-player (tracking).

And stop according to GLRT stopping rule.

Computational Complexity

Track-and-Stop: solves one “max-min” at each stage.

$$\operatorname{argmax}_{w \in \Delta_K} \inf_{\lambda \in \neg i^*(\hat{\mu}_t)} \sum_{k=1}^K w^k d(\hat{\mu}_t^k, \lambda^k).$$

AdaHedge + Best-response: solves one “min” at each stage.

$$\operatorname{argmin}_{\lambda \in \neg i_t} \sum_{k=1}^K w_t^k d(\hat{\mu}_t^k, \lambda^k).$$

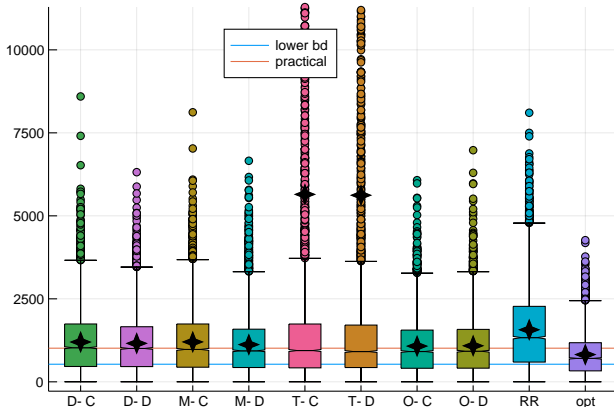
Examples

- Thresholding: closed form vs closed form.
- BAI: (line search)² vs line-search.
- Many Problems (sparse, lipschitz, unimodal): complicated? vs convex.

Results

For all $\mu \in \mathcal{M}$,

$$\mathbb{E}_{\mu} \tau_{\delta} \leq \frac{\log(1/\delta)}{\max \inf \sum_{k=1}^K w^k d(\mu^k, \lambda^k)} \left(1 + \mathcal{O} \left(\frac{1}{\sqrt{\log(1/\delta)}} \right) \right).$$



Variants

- Solve **max-max-min** at each stage \Rightarrow lowest sample complexity.
- Use a **learner for λ** \Rightarrow no tracking needed:
 - Follow the perturbed leader: always available but t samples at stage t ,
 - Easy if union of few simple convex regions.

Open problem

What if only few samples are available?

What if we want $\delta = 1/4$?

Conclusion

- Pure Exploration is a very broad setting.
- The game point of view is successful.
- Many other applications possible in bandits.
- The small confidence regime is still unclear.

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Thank you!